

Irwin-Hall distribution

Let $(U_i)_{i=1,\dots,m}$ be independent random variables uniformly distributed over $[-\frac{w}{2}, \frac{w}{2}]$. The distribution **Irwin-Hall** (m, w) is defined as the distribution of:

$$\sum_{i=1}^m U_i.$$

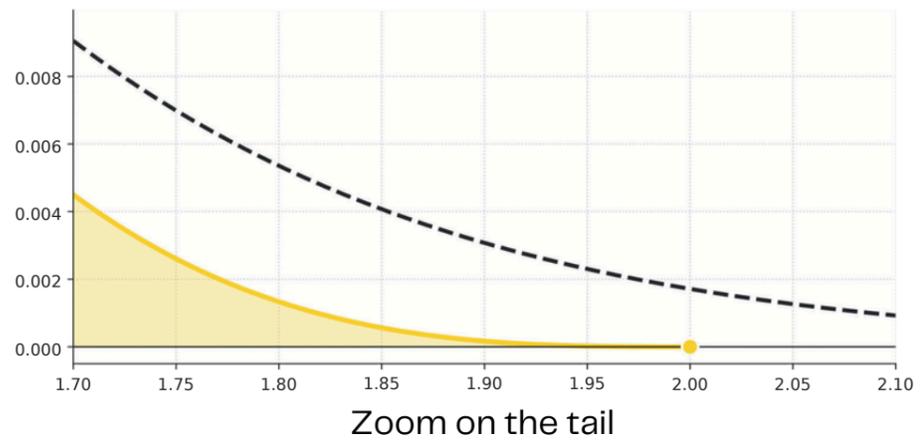
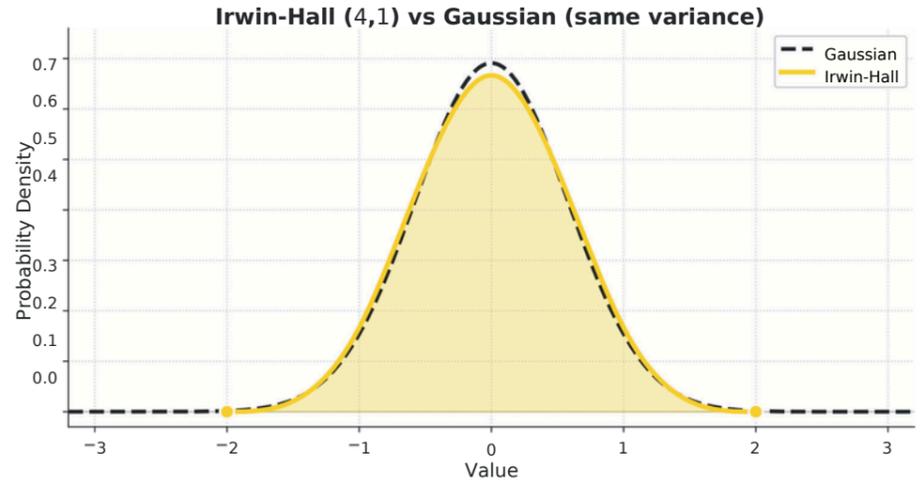
Modulus Switch noise: Irwin-Hall vs. Gaussian

Consider several variants of Modulus Switch (switching from modulus q' to q) for **LWE** ciphertexts with dimension n and binary secret key. The noise added by these operations is better modeled by an Irwin-Hall distribution than by a Gaussian distribution.

- Modulus Switch: **Irwin-Hall** $(\frac{n}{2}, \frac{1}{2q})$
- Mean Compensated [RDV25] **MS**: **Irwin-Hall** $(n, \frac{1}{4q})$
- Multi Bit [Zho+18] **MS**, grouping factor g :
Irwin-Hall $((1 - \frac{1}{2^g}) \frac{n}{g}, \frac{1}{2q})$

	Modulus Switch	Mean Comp. MS	Multi Bit MS
Importance Splitting	9.33e-23	3.28e-40	2.81e-43
(Upper Bound 3σ)	1.20e-22	4.34e-40	5.18e-43
Gaussian Prediction	2.20e-22	9.82e-40	3.70e-41
Irwin-Hall Prediction	1.05e-22	3.27e-40	2.20e-43

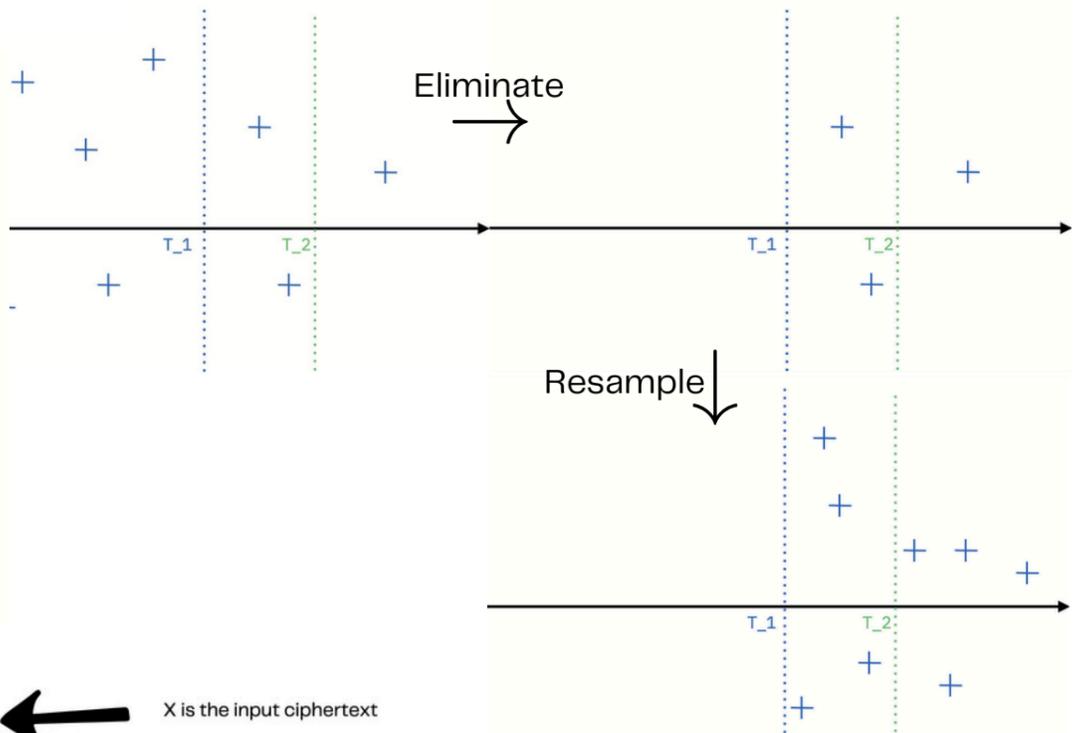
Table: Failure Probability: Experimental measures (Importance Splitting) vs. Theoretical Predictions.



Importance Splitting: the algorithm

- Fix a sequence of thresholds $T_1 < T_2 < \dots < T_l = T$
- Initialize the population $X_1^{(0)}, \dots, X_n^{(0)}$
- For $i = 1, \dots, l$ do
The population is $X_1^{(i-1)}, \dots, X_n^{(i-1)}$ and $S(X_j^{(i-1)}) > T_{i-1}$ for all j
 - Keep only the samples with score larger than T_i and count them
$$\{Y_1^{(i)}, \dots, Y_{n_i}^{(i)}\} := \{X_j^{(i-1)} \text{ such that } S(X_j^{(i-1)}) > T_i\}$$
 - Compute $p_i = \frac{n_i}{n}$
 - Resample a population of size n from the population $Y_1^{(i)}, \dots, Y_{n_i}^{(i)}$
$$X_1^{(i)}, \dots, X_n^{(i)}$$

such that $S(X_j^{(i)}) > T_i$ for all j
- Return $p = p_1 p_2 \dots p_l$



$S(X)$ is the noise after the modulus switch

Estimate Failure Probability of the whole atomic pattern

Input:

- Random Variable X
- A Score Function $X \rightarrow S(X)$

Output:

- Estimation of $P(S(X) > T)$
- Rare Events Generator

X is the input ciphertext

T is the correctness threshold
 $T = \Delta / 2$

[RDV25] Thomas de Ruijter, Jan-Pieter D'Anvers, and Ingrid Verbauwhede. *Don't be mean: Reducing Approximation Noise in TFHE through Mean Compensation*. Cryptology ePrint Archive, Paper 2025/809. 2025. URL: <https://eprint.iacr.org/2025/809>.

[Zho+18] Tanping Zhou, Xiaoyuan Yang, Longfei Liu, Wei Zhang, and Ningbo Li. "Faster Bootstrapping With Multiple Addends". In: *IEEE Access* 6 (2018), pp. 49868–49876. DOI: 10.1109/ACCESS.2018.2867655.

Bayes formula makes it work:

$$P(S > T_l) = P(S > T_1)P(S > T_2 | S > T_1) \dots P(S > T_l | S > T_{l-1})$$